

Social Choice Theory

▷ Preference Aggregation:

Suppose we know that every individual in society wants.
How can we aggregate these individual preferences into a social choice?

Social Choice Function: $\left\{ \begin{array}{l} \text{possible} \\ \text{choices} \end{array} \right\}^{\text{Individuals}} \rightarrow \left\{ \begin{array}{l} \text{possible} \\ \text{choices} \end{array} \right\}$

▷ Mechanism Design:

How can we design a voting procedure such that it is in everyone's best interest to actually reveal their true preferences?

▷ Power Indices:

How much power has an individual in a given voting procedure?

candidates,
choices
↓

▷ Most famous result: Arrow's Theorem

preference := linear ordering of a finite set C

social choice function:

$$f((\text{lin ord of } C)_{\text{voter}}) = \text{lin ord of } C$$

Arrow:

There is no "democratic" social choice function!

▷ huge literature of impossibility theorems.

▷ not the subject for today.

Spatial Voting

$x_0, x_1 \in X \leftarrow$ policy space

$x_0 \prec_i x_1 \iff$ voter $i \in V$ prefers x_1 to x_0

$x_0 \prec x_1 \iff$ the majority of voters prefers x_1 to x_0

What if $X \subset \mathbb{R}^n$?

Assumption:

\triangleright Each voter $i \in V$ has an ideal policy $p_i \in \mathbb{R}^n$.

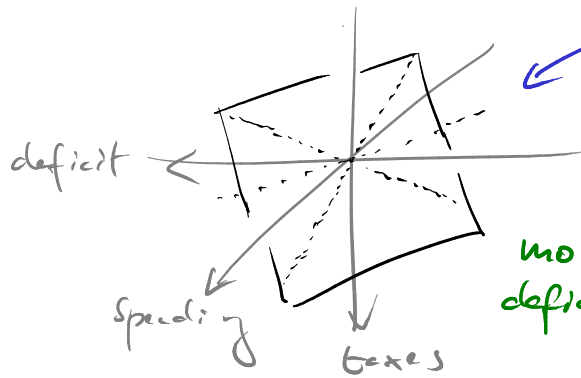
$\triangleright x_0 \prec_i x_1 \iff \|p_i - x_0\| < \|p_i - x_1\|$



Euclidean distance

Example: Budget \leftarrow (at least) 2-dimensional

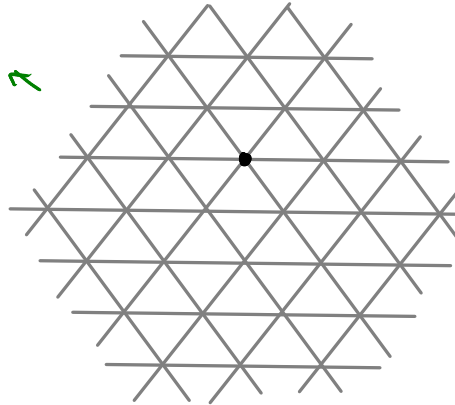
spending - taxes = deficit



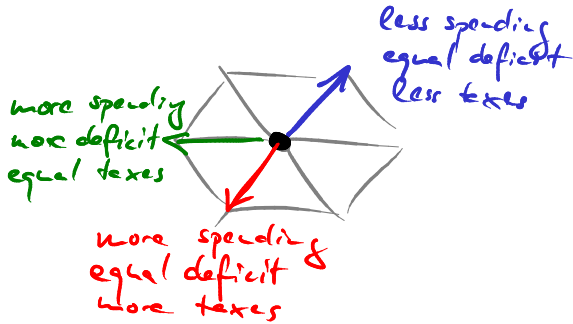
Budget decisions are points in this hyperplane

more deficit \leftarrow

\rightarrow less spending



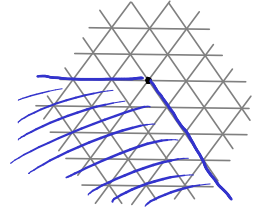
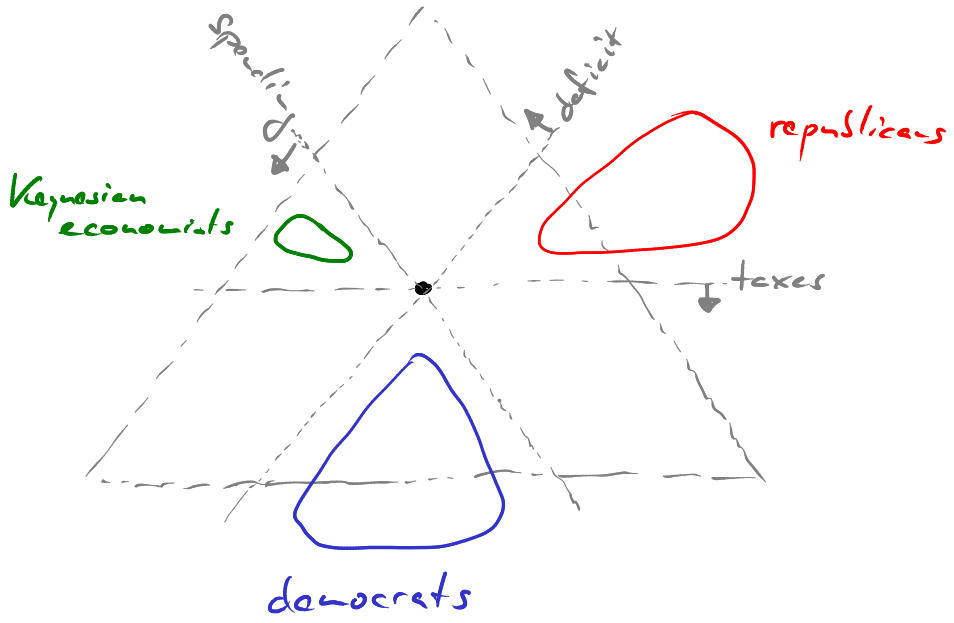
\downarrow
more taxes



more spending
more deficit
equal taxes

less spending
equal deficit
less taxes

more spending &
equal deficit &
more taxes



spending cut taxes
non-negative

Agenda Control

decide between two alternatives:

status quo



x_0

x_1



law policy
proposed by bill

McKelvey's Theorem

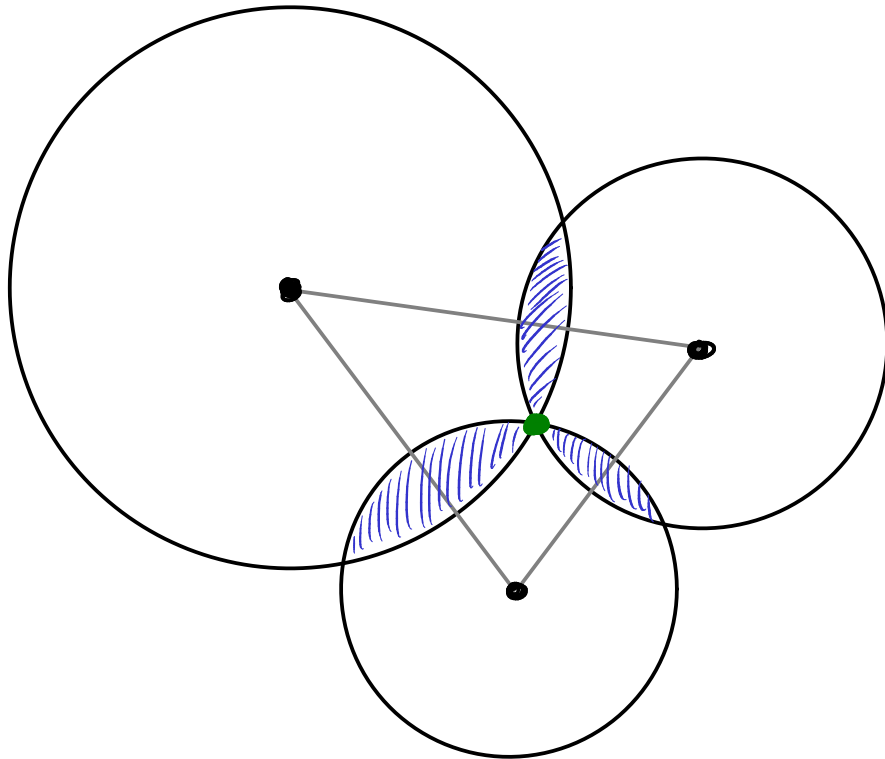
Given $x, y \in \mathbb{R}^n$. Then there exist $x_i \in \mathbb{R}^n$ such that

$$x = x_0 < \underbrace{x_1 < x_2}_{\text{second poll}} < x_3 < \dots < \underbrace{x_{n-1} < x_n}_{\text{n-th poll}} = y$$

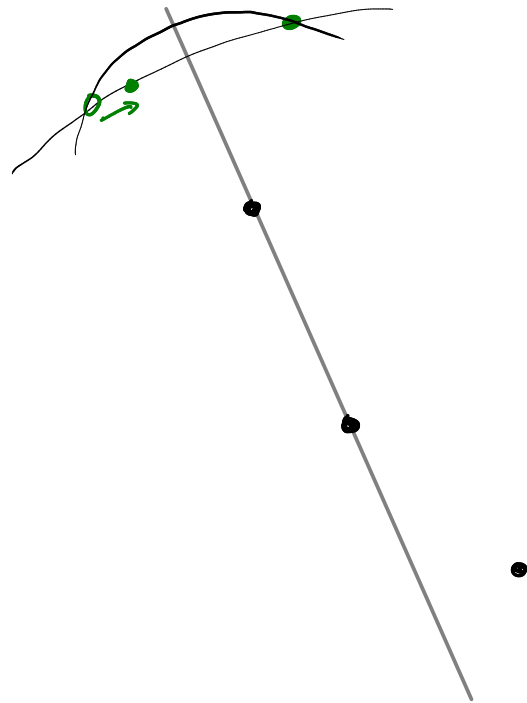
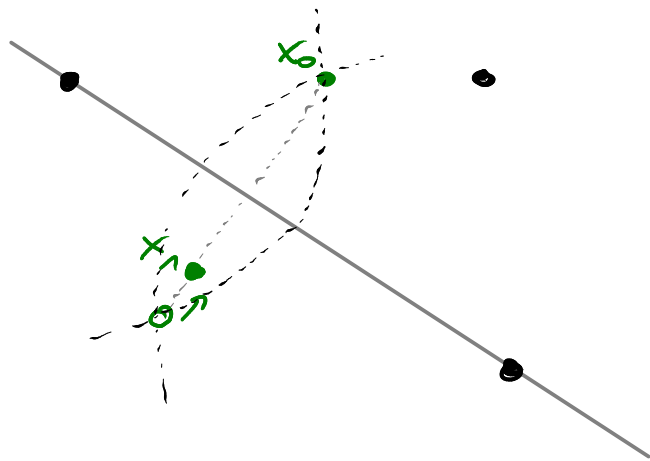
$\underbrace{x_0 < x_1}_{\text{first poll}}$ $\underbrace{x_2 < x_3}_{\text{third poll}}$

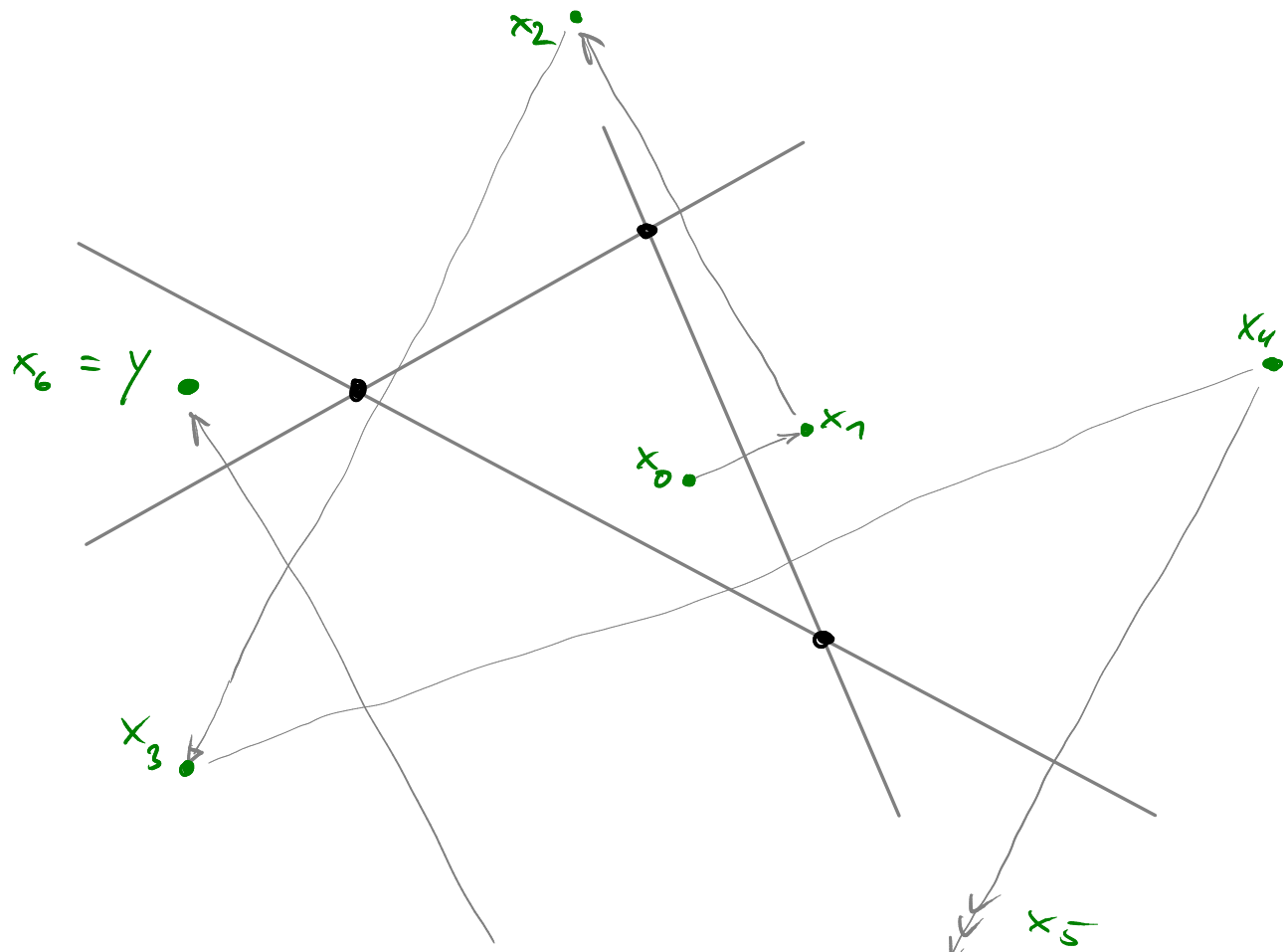
You can get voters to approve any bill, if you control the agenda.

Which x_1 defeat x_0 in a majority vote?



Reflection at one of these lines always works!





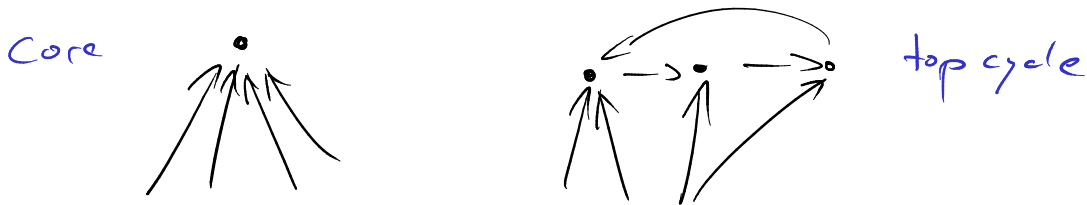
McKelvey's Theorem dimension ≥ 2 , # voters ≥ 3

If the core of the spatial voting game is empty,
its top cycle is \mathbb{R}^n .

Assume: # voters is odd.

Def: core = $\{x \mid y \prec x \ \forall y \in \mathbb{R}^n\}$.

top cycle = $\{x \mid \forall y \exists x_i : y \prec x_0 \prec \dots \prec x_n \prec x\}$

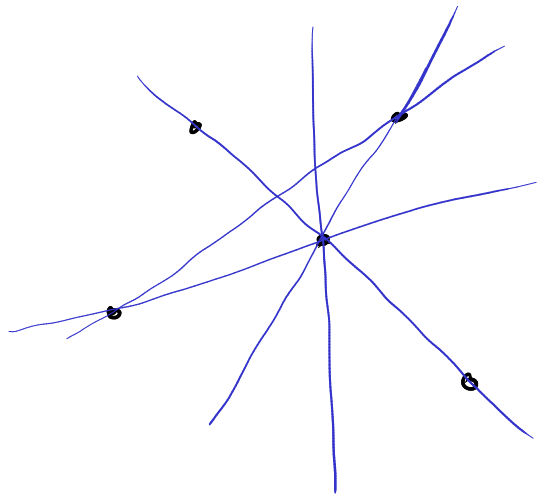


Thm: In general, the core is empty in $\dim \geq 2$.

Median Hyperplane Let $V = \{p_v \mid p_v \text{ ideal policy of voter } v\}$

For every normal a there exists b_a such that

$$|H_{a,b_a}^+ \cap V| \geq \frac{|V|}{2} \leq |H_{a,b_a}^- \cap V|$$

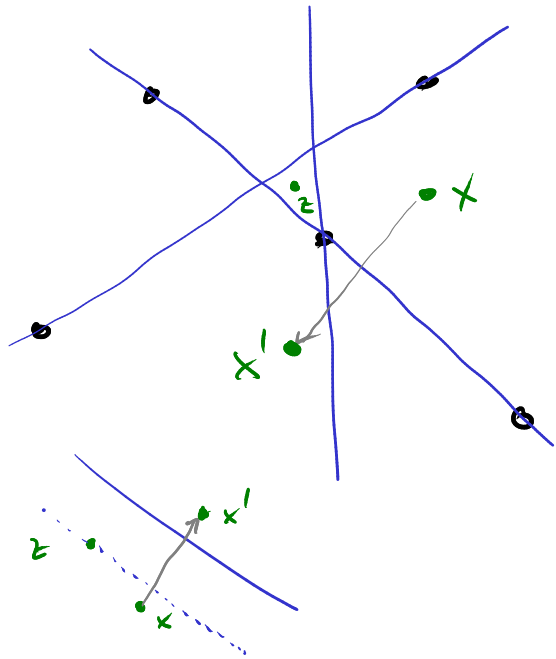


▷ # voters is odd $\Rightarrow S_a$ is unique

$$\text{core} = \bigcap_a H_{a,b_a}$$

┌ Let $x \notin H_{a,b_a}$. Let $\pi(x)$
be the orth. proj. of x onto
 H_{a,b_a} . Then $x \prec \pi(x)$. ┘

Proof Sketch for McKelvey's Thm



▷ Pick $d+1$ median hyperplanes

$$\bigcap_{i=1}^{d+1} H_{a_i, b_{a_i}}$$

▷ Pick a point z "inside the triangle",

▷ Show that $\forall x \in \mathbb{R}^d$ there ex i with

$x \notin H_{a_i, b_{a_i}}$ s.t. by moving

x "across" $H_{a_i, b_{a_i}}$ you get x' s.t.

$x \succ x'$ and $\|x - z\| < \|x' - z\|$.

▷ Construct $x_0 \prec x_1 \prec \dots \prec x_{n-1}$ that gets arbitrarily far away from z .

▷ If x_{n-1} is "far away enough", going to y will be an improvement for everybody. \square

Generalizations:

If $X \subset \mathbb{R}^n$ is connected, and voters have a continuous utility function, then, in general, the top cycle of the majority rule is always going to be all of X .

"Regardless of other voter's preferences, any one voter with complete information of the other voter's preferences, control of the agenda, and the ability to cast his own vote as he chooses can always construct majority paths to get anywhere in space."

Strong Point

$$\text{defeat}(x) = \text{Vol} \{ y \mid y \succ x \}$$

$x = \underset{x}{\text{argmin}} \text{defeat}(x)$ is called the **strong point**

Then x exists and is unique

Then in dimension 2:

$x = \sum \lambda_i v_i$ convex combination of voter's ideal policies v_i .

$\lambda_i = \underline{\text{spatial power index}}$ of voter i

Roughly: centrist voters have more power.